

bipartite
trees

S YSM 6302

CLASS 3

Bipartite Network

$$G(W, E) =$$

↑ nodes ↑ edges

$$e \in W \times W$$

$$B(U, V, E)$$

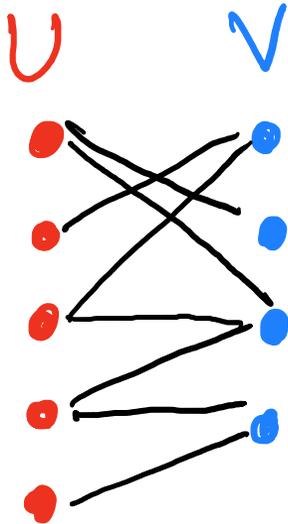
↑ nodes ↑ edges

Nodes can be partitioned into two types: "U" and "V" (these are disjoint)

$$W = U \cup V \quad e \in U \times V$$

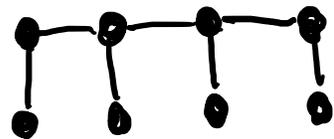
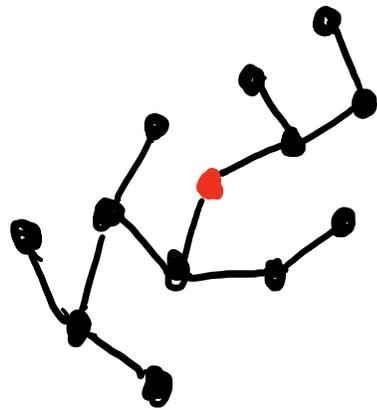
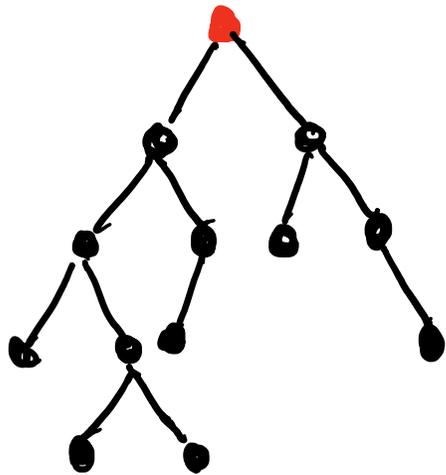
and

$$U \cap V = \emptyset$$



Tree

"rooted"



A tree is a connected, undirected network with no cycles.



↑ every vertex can be reached from every other vertex by a path

No cycles \Rightarrow there is only one path between any two nodes

Forest is a network made up of a collection of trees

↑ (disjoint union) \leftarrow they don't overlap

\Rightarrow there is at most one path between pairs of nodes (there may be no path)



A (finite) connected graph G with n nodes is a tree if and only if it has $n-1$ edges.

⇒ Suppose G is a tree

By induction

$n=1$

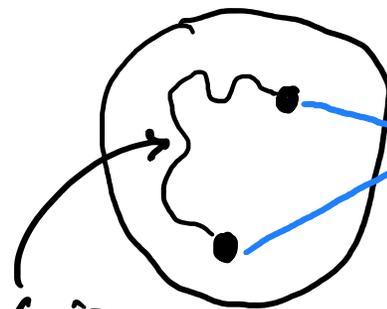


↑ this is a (trivial) tree
1 node, 0 edges

Suppose Theorem is true for n

∴ G has n nodes
 $n-1$ edges

To make a graph with $n+1$ nodes, need to add one node and some # of edges



new node

0 edges ⇒ disconnected

2+ edges ⇒ cycle

⇒ 1 edge!

Since G is a tree always one path b/t pairs of nodes

A (finite) connected graph G with n nodes is a tree if and only if it has $n-1$ edges.



⇐ Suppose G has $n-1$ edges

→ if it is a tree, we are done

→ if it is not a tree, it must contain a cycle

↳ so delete an edge in the cycle & repeat until a tree remains

↳ Call the new graph \tilde{G} , which has $n-1$ edges by previous slide

↳ So we started with G with $n-1$ edges, then removed

edges until we got \tilde{G} with $n-1$ edges ... Contradiction! G is a tree!

Directed Trees

A directed graph whose underlying graph is a tree. ↙ remove orientation

Not quite the same as an acyclic graph (the lab has you think about this)

directed acyclic graphs $\stackrel{?}{\subseteq}$ directed trees

OR

directed trees $\stackrel{?}{\subseteq}$ directed acyclic graphs

